Distributed On-Line Coordination for Multi-Robot Patrolling

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Motivation

RoCoCo **Cognitive Cooperating Robots** Lab. Research on Multi-robot systems



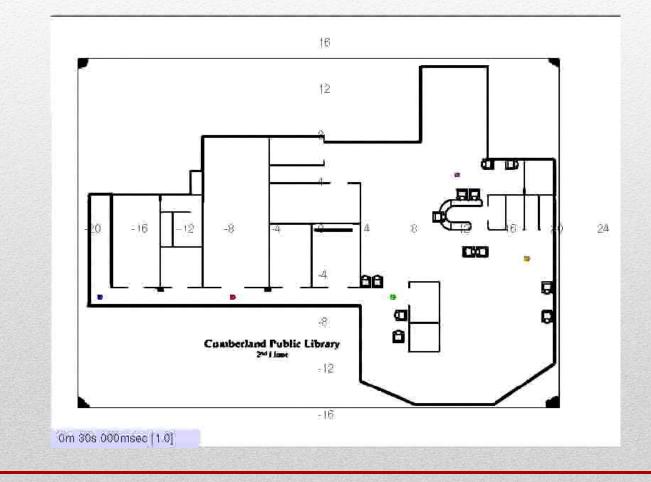
Motivation

Research projects on Multi-robot surveillance (mobile sensor network)

- Bee-SAFE Multi-robot mapping and exploration, (2011-2013), Sistemi Software Integrati, SELEX-SI
- Smart Monitoring of Complex Public Scenes, (2011-2013), Dept. of Homeland Security (DHS), USA
- Situation Awareness, Iniziativa Software 2, (2010-2012), SESM, SELEX-SI
- Multi-Robot Teams for Environmental Monitoring, (2009-2011), Dept. of Homeland Security (DHS), USA
- SAMAS Adaptive multi-robot services and its applications to demining, (2009- 2011), Space Software Italia.

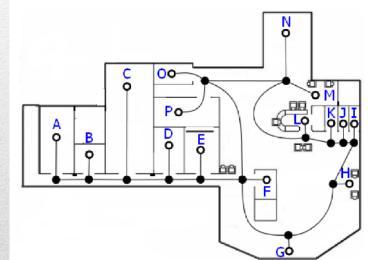
Multi-Robot Patrolling

Continuous monitoring of a (known) environment



MRP state of the art

- Graph based representation
- Performance metrics based on idleness and frequency
- <u>Perfect perception, action and</u> <u>communication</u>
- Optimal off-line centralized solutions
 - (e.g., derived using graph theory [Portugal and Rocha 11])



MRP our approach

- Graph based representation
- Performance metrics based on idleness and frequency
- Possible failures, robot interferences
- <u>(Suboptimal) on-line distributed</u> <u>solutions</u>
 - Based on Dynamic Task Assignment
- <u>Non-adversarial environments</u>
 - Visit places as often as possible



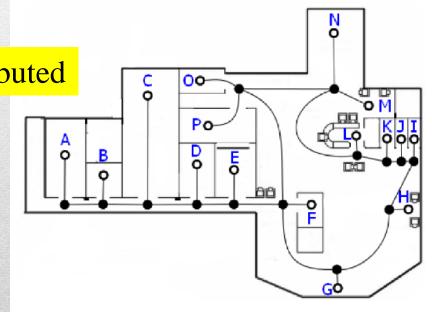
MRP Formulation

Patrol Graph $PG = \langle P, E, c \rangle$

P: Poses, E: Edges c: Travel cost

Objective On-line and distributed Find paths for each robot to minimize performance metrics

Idleness: $I^{p}(t)$ time between two consecutive visits of the same node



Dynamic Task Assignment for MRP

Assign robot **i** to task (location) $\mathbf{j} \rightarrow a_{ij} = 1$

Objective:

$$\arg\max_{\mathbf{A}} \sum_{i=1}^{|\mathcal{H}|} \sum_{j=1}^{|\mathcal{I}|} v_{ij} a_{ij}$$

|D| |D|

Reward of allocating **i** to **j**: $v_{ij} = U(r_i, p_j, t)$

$$U(r_i, p_j, t) = \theta_1 I^{p_j}(t) + \theta_2 T c(r_i, p_j, t)$$

Subject to:

- one robot per task \rightarrow avoid conflicts
- one task per robot \rightarrow avoid planning ahead



DTA Solutions

We aim at a completely **decentralized** solution:

- Each robot maintains $I^p(t)$ for all nodes
- Robots share info on idleness and merge such info
 - Merge: always keep lowest value for each room
- General solution scheme based on events
 - Robot reaches a visit location
 - Robot receives a message
- Two solutions
 - DTA-Greedy (always go to best location and deconflict)
 - DTAP (Sequential Single Item auctions and Partitioning)

DTA-Greedy

<u>When a robot reaches a location:</u> choose next <u>best</u> location (p*) set idleness of p* to 0 and broadcast <u>When receiving a message</u> (idleness) merge info on idleness revise destination



$$U(r_k, p', t) = \theta_1 I^{p_j}(t) + \theta_2 Tc(r_i, p_j, t) + \theta_3 ps(r_k, p')$$

Priority for location p'

- Defined off-line, easy to set on simple maps
- Balance work-load: e.g., low value for shared nodes



DTA with SSI and Partitioning

When a robot reaches a location:

- identify next <u>best</u> location (p*)
- announce \underline{bid} for target p^* (b)
- collect all bids
- if b is not the lowest, select next best location
- else consider robot responsible for p*, goto p*
- When receiving a message [idleness, bid, announce]
 - [idleness] merge info on idleness
 - [bid] updates bids
 - [announce] compute and send bid for target

DTAP bids

Bids based on utility function $U(r_i, p_j, t) = \theta_1 I^{p_j}(t) + \theta_2 T c(r_i, p_j, t)$

Cost based on a variable set of tasks assigned to a robot

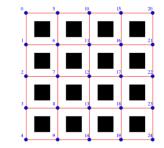
Variable set of tasks determine a **partition** of the graph

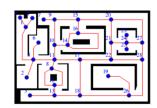
Difference with standard auctions:

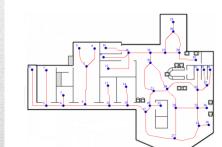
- no auctioneer
- tasks are not created but selected
- robot takes a task if no one else is better
- time-out when collecting bids (failures)

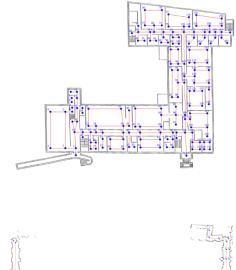
Experimental Set-up

- Realistic simulator (ROS & Stage)
- Six maps including real maps of our Dept.
- Many algorithms compared
- New performance representation and analysis











Performance metrics

Idleness: time between two consecutive visits of the same node by any robot in the team

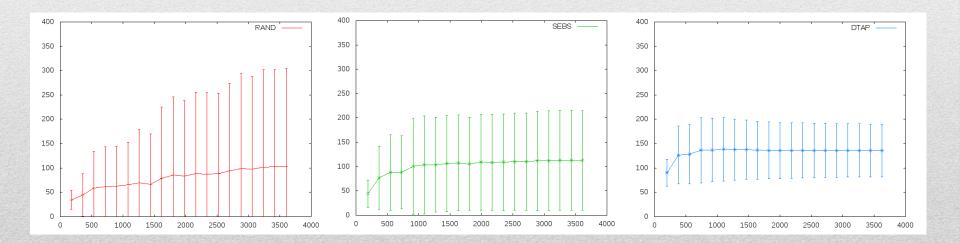
1000 800 800 Idleness (s) 400 200 0 0 500 1000 1500 2000 2500 3000 3500 Time (s)

Exp. ES6A: 4-3-4 robots cumberland_SEBS



Performance metrics

Idleness avg alone is not a good performance metric Idleness avg and stddev are independent metrics

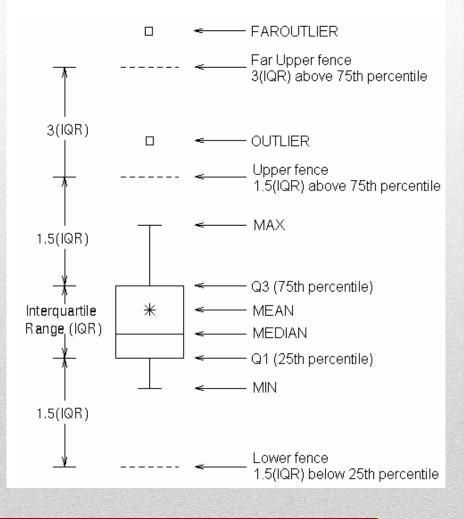




Performance metrics

Box-plot representation:

Graphical representation for distribution of idleness values during an experiment

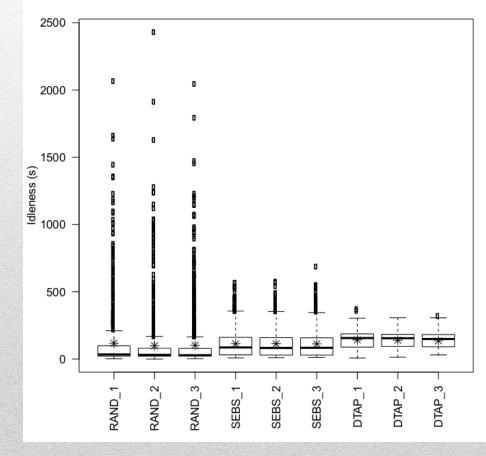




Results with simulator

Boxplots

- average
- stddev
- maximum
- outliers

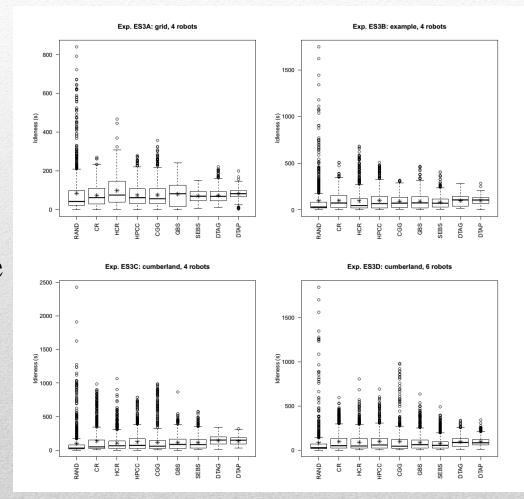


Exp. ES1: cumberland, 4 robots

Results with simulator

Comparison

- 9 algorithms (off-line vs. on-line)
- Best performance for on-line alg. minimize avg/stddev idl. outliers

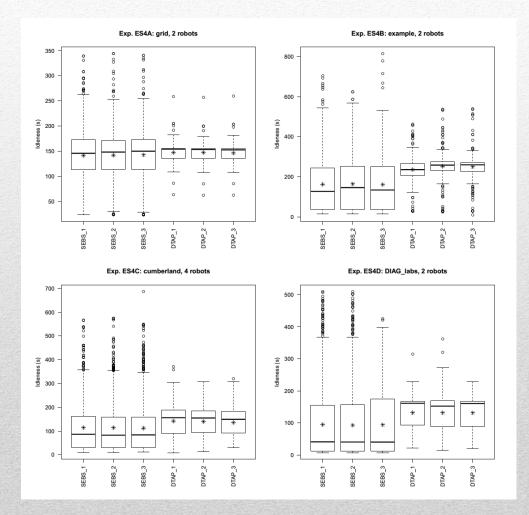




Results with simulator

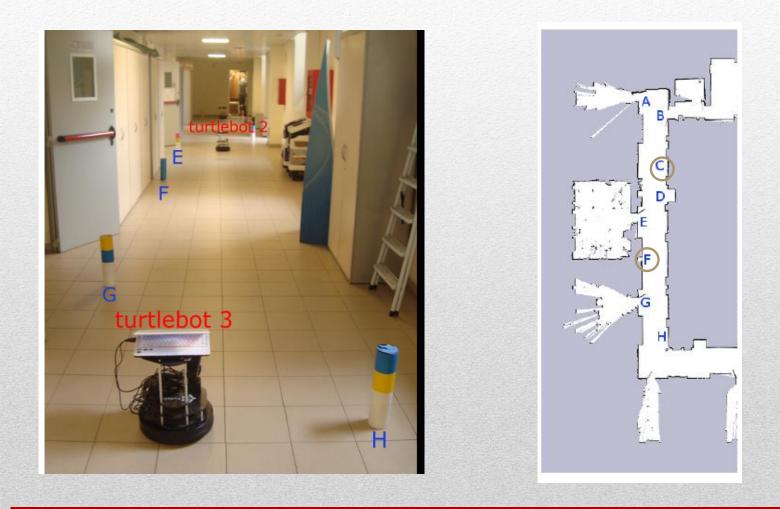
Comparison - 2 algorithms SEBS/DTAP

- **DTAP** higher avg lower stddev few outliers





Tests with robots



Tests with robots



Conclusions

Conclusions

Online DTA for practical approaches to MRP

 patrolling_sim available (standard benchmark for MRP) <u>http://wiki.ros.org/patrolling_sim</u>

Future Work

- Apply novel DTA approaches
 - improved robustness to comm. delays, failures
- Including learning techniques for adaptive behaviors